

# Game Theory - Week 1

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# Overview

- Course Overview
- What is a game?
- Normal Form Games
- Examples of Normal Form Games
- Constant Sum Games
- Games of Cooperation
- Nash Equilibrium

# Course Objectives

- Rigorous introduction to game theory and its applications to economics, political science, computer science, and biology.
- To give students a thorough understanding of how such problems are solved, and some experience in solving them
- To give students the background required to use the methods in their own research work or applications
- For more information visit course webpage <https://b2n.ir/r00714>

## Key Ingredients of Defining Game

- Who are the decision makers?
  - People? Governments? Companies? Somebody employed by a Company?
- What can the players do?
  - Enter a bid in an auction? Decide whether to end a strike? Decide when to sell a stock? Decide how to vote?
- What motivates players?
  - Do they care about some profit? Do they care about other players?

## Standard Representations of a Game

- 1 Normal form(Matrix form, Strategic form): Players take their actions simultaneously and payoffs are functions of players actions.
- 2 Extensive Form: Include timing of moves. Players move sequentially, represented as tree. Keeps track of what each player knows when he or she makes each decision

# Normal Form Games

## Definition (Normal Form Game)

A finite  $n$ -person normal form game is a triple  $\langle N, A, u \rangle$ :  
 $N = \{1, \dots, n\}$  a finite set representing players.  $A = A_1 \times \dots \times A_n$  is action profiles of players where each  $A_i$  present actions of player  $i$  and  $u = (u_1, \dots, u_n)$  is a profile of utility functions where each  $u_i$  is utility function of player  $i$ .

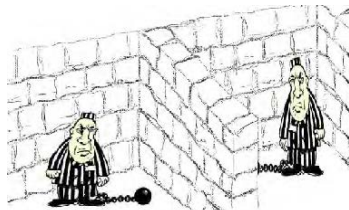
# Self-interested Agents

- What does it mean to say that an agent is self-interested?
  - Not that they want to harm others or only care about themselves
  - Only that the agent has its own description of states of the world that it likes, and acts based on this description
- Each such agent has a utility function
  - Quantifies degree of preference across alternatives
  - Explains the impact of uncertainty
  - Decision-theoretic rationality: act to maximize expected utility

# Prisoner's Dilemma

Each player have two option of Deny or Confess.

$$u_1(D, C) < u_1(C, C) < u_1(D, D) < u_1(C, D)$$



	D	C
D	(-1,-1)	(-10,0)
C	(0,-10)	(-8,-8)



# A Large Collective Action Game

- Players:  $N = \{1, \dots, 10000000\}$
- Action set of each player:  $A_i = \{Revolt, Not\}$
- Utility function of each player
  - $u_i(a) = 1$  if  $\#\{j | a_j = Revolt\} \geq 2000000$
  - $u_i(a) = -1$  if  $\#\{j | a_j = Revolt\} < 2000000$  and  $a_i = Revolt$
  - $u_i(a) = 0$  if  $\#\{j | a_j = Revolt\} < 2000000$  and  $a_i = Not$

# Constant Sum Games

Constant sum games also known as games of pure competition.

- Players have exactly opposed interests
- There must be precisely two players (otherwise they can't have exactly opposed interests)
- For all action profiles  $a \in A : u_1(a) + u_2(a) = c$  for some constant  $c$ . When  $c = 0$  game is called a zero sum game (e.g. pick a hand game).
- Thus, we only need to store a utility function for one player. In a sense, we only have to think about one player's interests

# Matching Pennies

One player want to match while other wants to mismatch

	H	T
H	(1,-1)	(-1,1)
T	(-1,1)	(1,-1)

## Pick a Hand

Hider(player I) may hide a coin in his left hand or two coins in his right hand. Chooser(player II) choose a hand and will get the coins in that hand.



	L	R
L	1	0
R	0	2

# Games of Cooperation

- Players have exactly the same interests
- No conflict: all players want the same thing
- $\forall a \in A : u_i(a) = u_j(a)$
- We often write such games with a single payoff per cell
- Why are such games “noncooperative”?

# Coordination Game

- Which side of the road you should drive on?

	Left	Right
Left	(1,1)	(0,0)
Right	(0,0)	(1,1)

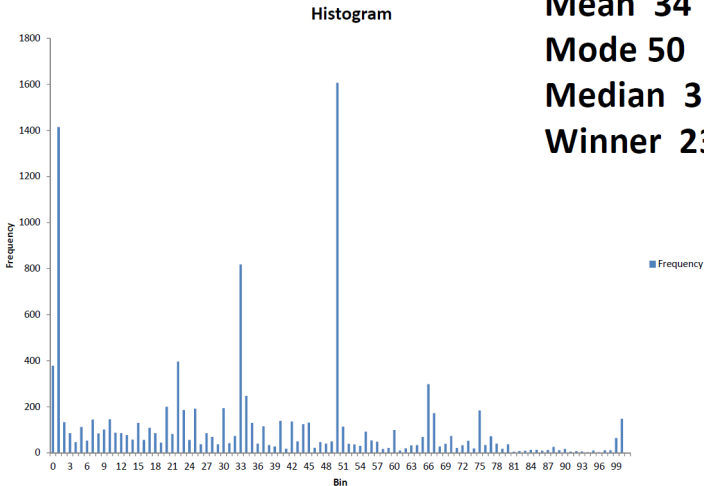
- Other examples: Battle of sexes, Stag hunt

## Keynes Beauty Contest Game

- Each player name an integer between 1 and 100
- The player who names the integer closest to two thirds of the average integer wins a prize, the other get nothing
- Ties are broken uniformly at random

Online course: more than 10000 players:

**2012 GTOC**  
**Mean 34**  
**Mode 50**  
**Median 33**  
**Winner 23**





# Best Response

- If you knew what everyone else was going to do, it would be easy to pick your own action
- let  $a_{-i} = (a_1, \dots, a_{i-1}, a_{i+1}, \dots, a_n)$ . And  $a = (a_{-i}, a_i)$

## Definition (Best Response)

$a_i^* \in BR(a_{-i})$  iff  $\forall a_i \in A_i : u_i(a_{-i}, a_i^*) \geq u_i(a_{-i}, a_i)$

# Nash Equilibrium

- Really, no agent knows what the others will do
- What can we say about which actions will occur ?
- Idea: look for stable action profiles

## Definition (Nash Equilibrium)

$a = (a_1, \dots, a_n) \in A$  is a (pure)Nash equilibrium iff  
 $\forall i: a_i \in BR(a_{-i})$

## Nash Equilibrium in Examples

	Left	Right
Left	(1,1)	(0,0)
Right	(0,0)	(1,1)

	H	T
H	(1,-1)	(-1,1)
T	(-1,1)	(1,-1)

	B	F
B	(2,1)	(0,0)
F	(0,0)	(1,2)

	D	C
D	(-1,-1)	(-4,0)
C	(0,-4)	(-3,-3)

# Strategies

From now we will call actions of each player (i.e. elements of  $A_i$ ) a pure strategy. We will face mixed strategies later in the course.

# Domination

- Let  $s_i$  and  $s'_i$  be two strategies for player  $i$  and let  $S_{-i}$  be the set of all possible pure strategy profiles for the other players

## Definition

$s_i$  strictly dominates  $s'_i$  if  $\forall s_{-i} \in S_{-i} : u_i(s_i, s_{-i}) > u_i(s'_i, s_{-i})$

## Definition

$s_i$  weakly dominates  $s'_i$  if  $\forall s_{-i} \in S_{-i} : u_i(s_i, s_{-i}) \geq u_i(s'_i, s_{-i})$

## Equilibria and Dominance

- If one strategy dominates all others, we say it is dominant.
- A strategy profile consisting of dominant strategies for every player must be a Nash equilibrium
- An equilibrium in strictly dominant strategies must be unique.

	D	C
D	$(-1,-1)$	$(-4,0)$
C	$(0,-4)$	$(-3,-3)$

# Pareto Optimality

- Idea: sometimes, one outcome  $o$  is at least as good for every agent as another outcome  $o'$ , and there is some agent who strictly prefers  $o$  to  $o'$
- In this case, it seems reasonable to say that  $o$  is better than  $o'$
- We say that  $o$  Pareto-dominates  $o'$

## Definition (Pareto Optimality)

An outcome  $o^*$  is Pareto-optimal if there is no other outcome that Pareto-dominates it.

- Can a game have more than one Pareto-optimal outcome?
- Does every game have at least one Pareto-optimal outcome?